

Patch Locale of a Spectral Locale in Univalent Type Theory

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What is a locale?

Notion of space characterised solely by its **frame of opens**.

Spectral and Stone locales

Spectral

Stone

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A locale in which the compact opens form a basis that is closed under finite meets.

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A compact locale in which the clopens form a basis.

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Space of ideals of **distributive lattice**.

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A **compact** locale in which **the clopens form a basis**.

Space of ideals of a **Boolean algebra**.

Morphisms are **continuous maps**.

Every Stone locale is spectral.

Spectral and Stone locales

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A **compact** locale in which **the clopens form a basis**.

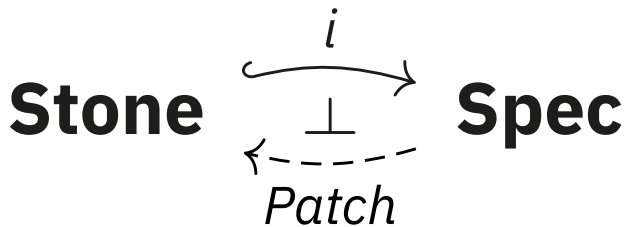
Space of ideals of a **Boolean algebra**.

Morphisms are **continuous maps**.

Every Stone locale is spectral.

Patch universally transforms a spectral locale into a Stone one.

Patch as a coreflector



Some examples of patch

Spectral locale in consideration

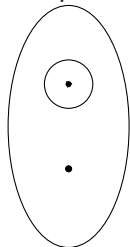
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Sierpiński space (Ω)



Scott topology of a (Scott) domain

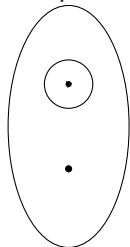
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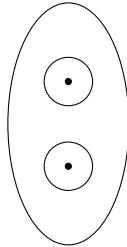
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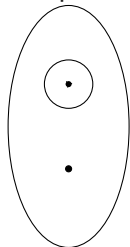
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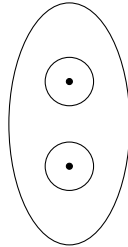
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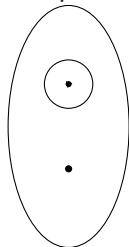


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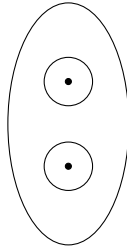
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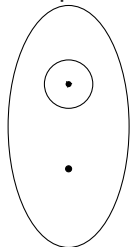
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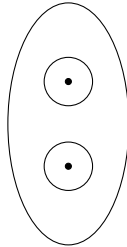
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\mathbb{N}_{∞}

Goal

Implement the patch locale in
univalent type theory
predicatively i.e. without using
propositional resizing axioms.

Type-theoretical preliminaries and notation

Definition (Family)

The type of \mathcal{W} -families over a type $A : \mathcal{U}$ is $\text{Fam}_{\mathcal{W}}(A) \equiv \Sigma_{I:\mathcal{W}} (I \rightarrow A)$.

Type-theoretical preliminaries and notation

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A **subfamily** of a \mathcal{W} -family (I, α) is a family $(J, \alpha \circ \beta)$ where (J, β) is itself a \mathcal{W} -family on I .

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Definition (Propositional resizing)

The **propositional resizing axiom** from \mathcal{U} to \mathcal{V} says

$$\prod_{P:\mathcal{U}} \text{is-prop}(P) \rightarrow \text{is-}\mathcal{V}\text{-small}(P).$$

Frames in type theory

Definition (Frame)

A $(\mathcal{U}, \mathcal{V}, \mathcal{W})$ -frame consists of

- ▶ a type $A : \mathcal{U}$,
- ▶ a partial order $- \leq - : A \rightarrow A \rightarrow \mathbf{hProp}_{\mathcal{V}}$,
- ▶ a top element $\top : A$,
- ▶ a binary meet operation $- \wedge - : A \rightarrow A \rightarrow A$,
- ▶ a join operation $\bigvee _ : \mathbf{Fam}_{\mathcal{W}}(A) \rightarrow A$,
- ▶ satisfying distributivity i.e. $x \wedge \bigvee_{i:I} y_i = \bigvee_{i:I} x \wedge y_i$ for every $x : A$ and family $\{y_i\}_{i:I}$ in A .

Some notation

A **frame homomorphism** is a function preserving finite meets and arbitrary joins.

The category of frames and their homomorphisms is denoted **Frm**; its opposite is denoted **Loc**.

- ▶ Morphisms of **Loc** are called **continuous maps**.

The frame corresponding to a locale X is denoted $\mathcal{O}(X)$.

We work in the spatial direction:

- ▶ X, Y, Z, \dots range over locales;
- ▶ $f, g : X \rightarrow Y$ range over continuous maps;
- ▶ $U, V, W, \dots : \mathcal{O}(X)$ range over opens; and
- ▶ $f^* : \mathcal{O}(Y) \rightarrow \mathcal{O}(X)$ denotes the frame homomorphism corresponding to a continuous map $f : X \rightarrow Y$ of locales.

Notions of size for locales

The generality of $(\mathcal{U}, \mathcal{V}, \mathcal{W})$ -locales is **almost never** needed.

Definition

A **small locale** is a $(\mathcal{U}, \mathcal{U}, \mathcal{U})$ -locale.

A **large** and **locally small** locale is a $(\mathcal{U}^+, \mathcal{U}, \mathcal{U})$ -locale.

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Theorem (de Jong and Escardó [6] [5])

The existence of a nontrivial small locale implies a form of propositional resizing.

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Most locales that arise in practice in our setting are
large and locally small.

Bases for locales

Definition (Basis)

A \mathcal{W} -family $\{B_i\}_{i:I}$ over a $(\mathcal{U}, \mathcal{V}, \mathcal{W})$ -locale X is said to **form a basis** for X if

for any $U : \mathcal{O}(X)$, there is a *subfamily* $\{B_l\}_{l \in L}$ of $\{B_i\}_{i:I}$ such that

$$U = \bigvee_{l \in L} B_l.$$

Note that a locale basis can be assumed WLOG to be directed.

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In our work, we focus on:

large and locally small frames with small bases.

A description of Patch

A nucleus on frame L is an endofunction $j : |L| \rightarrow |L|$ that is inflationary, idempotent, and preserves binary meets.

A nucleus is called **Scott-continuous** if it preserves joins of directed families.

Patch is the frame of **Scott-continuous nuclei**.

Previous work [3] exploited the fact that Patch is a subframe of the frame of all nuclei.

Joins in the frame of *all* nuclei (1)

Consider a locale X and let $j, k : \mathcal{O}(X) \rightarrow \mathcal{O}(X)$ be two nuclei.

Ordering: $j \preceq k \equiv \prod_{U:\mathcal{O}(X)} j(U) \leq k(U)$

Top element: $\mathbf{1} \equiv U \mapsto \mathbf{1}_X$.

Binary meets: $j \wedge k \equiv U \mapsto j(U) \wedge k(U)$.

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Binary meets: $j \wedge k \equiv U \mapsto j(U) \wedge k(U)$.

Unfortunately, the **pointwise join** fails to be idempotent in general.

Joins in the frame of *all* nuclei (2)

It is possible to construct the joins in the frame of all nuclei **impredicatively**.

Previous constructions include those by

- ▶ Simmons [7],
- ▶ Banaschewski [1],
- ▶ Johnstone [4],
- ▶ Wilson [8], and
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All of these use some form of **impredicativity**.
Moreover, we are not aware of any predicative construction of the frame of all nuclei.

The frame of Scott-continuous nuclei in type theory?

Question

Can this patch construction be shown to be the coreflector in the predicative setting of univalent type theory?

The frame of Scott-continuous nuclei in type theory?

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Our answer: **Yes**, if one considers **large** and **locally small** frames with **small bases**.

Joins in the frame of Scott-continuous nuclei (1)

Let $K := \{k_i\}_{i:I}$ be a family of nuclei.

Finite compositions of nuclei

Denote by K^* the family of **all finite compositions** of k_i 's:

$$K^* := (\text{List}(I), \beta)$$

where $\beta(i_0, \dots, i_{n-1}) := k_{i_{n-1}} \circ \dots \circ k_{i_0}$.

Proposition

The family K^* is always directed.

Joins in the frame of Scott-continuous nuclei (2)

When only the Scott-continuous nuclei are considered, however, this join exists **predicatively!**

Let $K \equiv \{k_i\}_{i:I}$ be a family of Scott-continuous nuclei on locale X and let $U : \mathcal{O}(X)$.

$$\left(\bigvee_{i:I} k_i\right) \left(\bigvee_{i:I} k_i\right) (U) \equiv$$

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Compactness

Definition (Compact open)

Let $U : \mathcal{O}(X)$ be an open of some locale X . The open U is called **compact** if

$$U \leq \bigvee_{i:I} V_i \rightarrow \exists k : I. U \leq V_k$$

for every *directed* family of opens $\{V_i\} : \text{Fam}_{\mathcal{W}}(\mathcal{O}(X))$.

Definition (Compact locale)

A locale X is called **compact** if its top open is compact.

The type-theoretical construction of Patch

Recall the definition of a spectral locale:

a locale in which the compact opens form a basis closed under finite meets.

One can predicatively construct the patch locale on an arbitrary locale.

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The need for a small basis arises in multiple places:

- ▶ to ensure that Patch is locally small,
- ▶ to ensure the existence of open nuclei (to be explained later).

Ordering on nuclei – *size matters*

Let

- ▶ X be a **large** and **locally small** locale, and
- ▶ j and k be two Scott-continuous nuclei on X .

Define $j \preceq k \equiv \prod_{U:\mathcal{O}(X)} j(U) \leq k(U)$.

Problem: $j \preceq k$ lives in universe \mathcal{U}^+ .

This means $\text{Patch}(X)$ is a $(\mathcal{U}^+, \mathcal{U}^+, \mathcal{U})$ -locale i.e. it is *not* **locally small**.

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Definition

$$j \preceq_s k \equiv \prod_{C:\mathcal{K}(X)} j(C) \leq k(C)$$


Closed and open nuclei

We embed the opens of X into $\text{Patch}(X)$ using the **closed** and **open** nuclei.

- ▶ **Open nucleus** induced by U :
 - ▶ Represents the **open subspace** corresponding to U ;
 - ▶ Defined as $\neg'U' :\equiv V \mapsto U \Rightarrow V$.
- ▶ **Closed nucleus** induced by U ,
 - ▶ Represents the **closed subspace** corresponding to the complement of U ;
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Heyting implication is defined using the Adjoint Functor Theorem which amounts to the definition:

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To solve these problems, we start with a spectral locale X assumed to

- ▶ be **large** and **locally small**;
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We make exactly the same assumptions for **zero-dimensionality**, involved in the definition of a Stone locale.

Spectrality revisited (2)

Proposition

In any spectral locale with small basis, any compact open falls in the basis.

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Proof

- ▶ Let $C : \mathcal{O}(X)$ be a compact open.
- ▶ $C = \bigvee_{l \in L} B_l$ for some basic covering family $\{B_l\}_{l \in L}$.
- ▶ By compactness, there exists an index $k \in L$ such that $C \leq B_k$.
- ▶ Clearly, $B_k \leq C$ which means $B_k = C$.

Spectrality revisited (3)

We can do better! Let X be a spectral locale with small basis $\{B_i\}_{i:I}$.

Theorem

There is an isomorphism

$$(\sum_{U:\mathcal{O}(X)} \exists_{i:I} U = B_i) \simeq \sum_{U:\mathcal{O}(X)} \text{is-compact-open}(U).$$

Proposition

Assuming the existence of quotients, the type $\sum_{U:\mathcal{O}(X)} \exists_{i:I} U = B_i$ is \mathcal{U} -small.

Corollary

The type of compact opens of X is \mathcal{U} -small.

Patch is Stone

Theorem

Given a spectral locale A with small basis $\{B_i\}_{i:I}$, $\text{Patch}(A)$ is a Stone locale (with a small basis consisting of clopens).

Proof sketch

The family

$$\{ 'B_m' \wedge \neg 'B_n' \mid m, n : I \}$$

forms a basis for $\text{Patch}(A)$ and the covering subfamily for a given Scott-continuous nucleus $j : \mathcal{O}(X) \rightarrow \mathcal{O}(X)$ is given by

$$\{ 'B_m' \wedge \neg 'B_n' \mid B_m \leq j(B_n), m, n : I \}.$$

The universal property of Patch

Let A and X be, respectively, a spectral and a Stone locale with small bases. Denote by $\{B_i\}_{i:I}$ the small basis of A .

$$\begin{array}{ccc} X & & \\ \downarrow f & \searrow \bar{f} & \\ A & \xleftarrow{\varepsilon} & \text{Patch}(A) \end{array}$$

The universal property of Patch

Let A and X be, respectively, a spectral and a Stone locale with small bases. Denote by $\{B_i\}_{i:I}$ the small basis of A .

$$\begin{array}{ccc} X & & \\ \downarrow f & \searrow \bar{f} & \\ A & \xleftarrow{\varepsilon} & \text{Patch}(A) \end{array}$$

$$\bar{f}^* : \mathcal{O}(\text{Patch}(A)) \rightarrow \mathcal{O}(X)$$

$$\bar{f}^* := j \mapsto \bigvee_{m,n:I} \{f^*(B_m) \wedge \neg f^*(B_n) \mid B_m \leq j(B_n)\}$$

Summary

We set out to implement a predicative version of an impredicative construction from pointfree topology in univalent type theory.

Doing this predicatively turned out to involve surprising challenges.

We had to reformulate quite a few things in the theory itself to obtain a **type-theoretic understanding** of the construction in consideration.

I have almost completely formalised our work in the Agda proof assistant, as part of Escardó's `TypeTopology [0]` library.

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